

# Optimal Recovery and $n$ -Widths for Convex Classes of Functions

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We study the problem of optimal recovery in the case of a nonsymmetric convex class of functions. In particular we show that adaptive methods may be much better than nonadaptive methods. We define certain Gelfand-type widths that are useful for nonsymmetric classes and prove relations to optimal error bounds for adaptive and nonadaptive methods, respectively. © 1995 Academic Press, Inc.

## 1. INTRODUCTION

No approximation scheme can be good for every function  $f$ . We need some a priori information about  $f$  of the form  $f \in F$ . Usually one assumes that  $f$  is an element of a certain Banach space  $X$  and so might have certain smoothness properties. Then it is our task, for example, to find a good approximation of the linear operator  $S: X \rightarrow G$  such that

$$\left\| S(f) - \sum_{i=1}^n L_i(f) \cdot g_i \right\|_G \leq c_n \cdot \|f\|_X$$

holds with as small a  $c_n$  as possible. Here the  $L_i$  are linear functionals,  $L_i: X \rightarrow \mathbf{R}$ , for example, function values or Fourier coefficients.

Often, we perform a worst case analysis on a unit ball

$$F = \{f \in X \mid \|f\|_X \leq 1\}$$

which is convex and symmetric, i.e.,  $-f \in F$  if  $f \in F$ . This approach is the usual one in numerical analysis, at least if the solution operator is linear. Also, most of the known results on optimal recovery and closely related problems on  $n$ -widths usually are studied under the assumption that the set  $F$  of problem elements is convex and symmetric. In many cases, however, we have a different type of a priori information. We give some examples.

Sometimes we know that  $f$  is positive because, for example,  $f$  is a certain density function. In this case we should consider sets of the type

$$F = \{f \in X \mid \|f\|_X \leq 1, f \geq 0\}.$$

Observe that such a set is still convex, but not symmetric. In other cases we might know in advance that  $f$  is a monotone or convex function. This also leads us to study convex classes of functions that are nonsymmetric. The geometric information given by positivity, monotonicity, or convexity is very important in some cases. It often helps to find an effective numerical method, even if the problem is ill-posed without this information.

Therefore it is usually not a good idea to just ignore the additional information about  $f$ . However, it may seem that it is still enough to study symmetric and convex sets—at least modulo some minor details. Let us again consider the case where we want to approximate a linear operator  $S$  on  $F$ . By taking  $F - F$ , defined by

$$F - F = \{f_1 - f_2 \mid f_1, f_2 \in F\},$$

we clearly get a symmetric set and for each convex set  $F$  we get the error estimate

$$\inf_{S_n} \Delta_{\max}^F(S_n) \leq \inf_{S_n} \Delta_{\max}^{F-F}(S_n) \leq 4 \inf_{S_n} \Delta_{\max}^F(S_n). \quad (1.1)$$

Here the maximal error is defined in the usual way by

$$\Delta_{\max}^F(S_n) = \sup_{f \in F} \|S(f) - S_n(f)\|$$

and the infimum runs through all methods of the form

$$S_n(f) = \phi(L_1(f), \dots, L_n(f)) \quad (1.2)$$

with nonadaptively chosen linear functionals  $L_i$ ; see Proposition 2. In the symmetric case we know that such nonadaptive methods are almost optimal in the class of all adaptive methods that use  $n$  linear functionals.

So we know that optimal error bounds for  $F$  and for  $F - F$  differ at most by a factor of 4 in the case of nonadaptive methods and adaption does not help (up to a factor of 2) for  $F - F$ . Therefore we can get much better error bounds on  $F$  only if we allow adaptive methods. We will see later that for some linear problems  $S: X \rightarrow G$  and convex  $F \subset X$  adaptive methods actually are much better than nonadaptive ones. This also proves that an inequality such as (1.1) does not hold if we allow adaptive methods.

We mention some of the literature in this subject. In the linear theory, i.e., under the assumption that  $F$  is symmetric and convex, the close connection between optimal recovery and  $n$ -widths or  $s$ -numbers is well

known; see Mathé [18], Micchelli and Rivlin [19, 20], Novak [23], Pinkus [29], Traub and Woźniakowski [36], and Traub, Wasilkowski, and Woźniakowski [35]. Useful surveys on  $n$ -widths are Pietsch [27], Pinkus [28], and Tikhomirov [34].

In the nonsymmetric case not so much is known. Some of the known  $n$ -widths can also be defined in the nonsymmetric case, but there is no theory of diameters in connection with optimal recovery, in particular when adaptive methods are allowed. Some special problems, however, are studied in the literature. The problem of optimal numerical integration of monotone functions was studied by Kiefer [12] and Novak [24]. The knots  $t_i$  may be chosen adaptively, i.e., sequentially. Kiefer proved that the best method is given by the trapezoidal rule. Hence we have an affine and nonadaptive algorithm which is optimal. Observe that adaption does not help in this case. This is also known for arbitrary linear  $S: F \rightarrow \mathbf{R}$  in the case where  $F$  is convex and symmetric; see Bakhvalov [3].

In the present paper we study the question of whether adaption can help if  $F$  is only convex. Also, in some other papers linear problems (such as integration or optimal reconstruction in  $L_{\infty}$ -norm) have been studied for certain nonsymmetric convex classes of monotone or convex functions. We mention the papers of Braß [6], Glinkin [8, 9], Novak [25], Petras [26], and Sonnevend [30].

Different nonsymmetric extremal problems in approximation theory were investigated by Babenko [1, 2], Gal and Micchelli [7], Ioffe and Tikhomirov [10], Korneichuk [15], Magaril-Ilyayev and Osipenko [17], Sukharev [31], and Sun [32]. We are mainly interested in the following question, where the worst case setting is studied for linear problems: Can adaption help (much) on a convex class of functions? Much is known about linear problems

$$S: X \rightarrow Y,$$

when considered on a symmetric and convex set  $F \subset X$  in the worst case. A slight superiority of adaptive methods can be proven in some cases even if  $F$  is symmetric; see Kon and Novak [13, 14]. It is well known, however, that adaption cannot help much in that case. Although adaptive methods are widely used, most theoretical results show that adaption does not help under various conditions.

It is known, however, that there are examples of a convex and nonsymmetric set  $F$ , where adaption helps considerably; see Novak [25] and Section 4. In this paper we define certain new "Gelfand-type"  $n$ -widths that turn out to be important for the study of linear problems on convex domains. We study the connection between these  $n$ -widths and problems of optimal recovery.

We believe that it is important to calculate the  $n$ -widths for standard classes of nonsymmetric sets, for example, sets of the type

$$\left\{ f: [0, 1] \rightarrow \mathbf{R} \left\{ \sum_{i=0}^k \|f^{(i)}\|_p \leq 1 \right\} \right\} \cap \{f \in C^l([0, 1]) \mid f^{(l)} \geq 0\}.$$

This would be useful for the construction of efficient algorithms for many practical problems.

In Section 5 we study the case where only methods of the form

$$S_n^{(\text{ad})}(f) = \phi(f(t_1), \dots, f(t_n)),$$

with function values instead of general linear functionals, are admissible. In this case adaption can help even more. An artificial example with exponential improvement is presented. A recent example from Korneichuk [16] which also shows that adaption can help is given.

## 2. DIAMETERS FOR NONSYMMETRIC SETS

We want to know whether adaption can help for linear problems on a convex set of functions. We begin slightly more generally and first define certain diameters that are interesting in the case where  $F$  is not symmetric.

Let  $X$  be a Banach space over  $\mathbf{R}$  and let  $F \subset X$  be convex. We first assume that  $F$  is also symmetric, i.e.,  $f \in F$  implies  $-f \in F$ . The Kolmogorov  $n$ -width of  $F$  in  $X$  is given by

$$d_n(F) = \inf_{X_n} \sup_{f \in F} \inf_{g \in X_n} \|f - g\|, \quad (2.1)$$

where the left infimum is taken over all  $n$ -dimensional subspaces  $X_n$  of  $X$ . Similarly, the Gelfand  $n$ -width of  $F$  is given by

$$d^n(F) = \inf_{U_n} \sup_{f \in F \cap U_n} \|f\|, \quad (2.2)$$

where the infimum is taken over all closed subspaces  $U_n$  of  $X$  with codimension  $n$ . These numbers measure the "thickness" or "massivity" of  $F$ .

In the case of arbitrary (in particular: nonsymmetric) sets  $F \subset X$  these definitions seem to be inadequate. The widths should be translation-invariant; therefore the Kolmogorov  $n$ -width (for arbitrary  $F \subset X$ ) should be given by

$$d_n(F) = \inf_{X_n} \sup_{f \in F} \inf_{g \in X_n} \|f - g\|, \quad (2.3)$$

where  $X_n$  runs through all  $n$ -dimensional affine subspaces of  $X$ ; see Tikhomirov [33]. For a convex and symmetric set  $F \subset X$ , (2.2) can be rewritten as

$$d^n(F) = \frac{1}{2} \cdot \inf_{U_n} \text{diam}(F \cap U_n).$$

Here,  $\text{diam}(B)$  means the diameter of a set  $B$ , defined by

$$\text{diam}(B) = \sup_{f, g \in B} \|f - g\|.$$

It is interesting to note that this definition (for symmetric sets) can be extended to arbitrary sets in two different ways. Both of them are interesting—at least if we are thinking of applications in the field of optimal recovery. A “global” variant of the Gelfand width (for arbitrary  $F$ ) is given by

$$d_{\text{glob}}^n(F) = \frac{1}{2} \cdot \inf_{U_n} \sup_{f \in X} \text{diam}(F \cap (U_n + f)) \quad (2.4)$$

(a slightly different notion is defined in Ioffe and Tikhomirov [10]), while a “local” variant is given by

$$d_{\text{loc}}^n(F) = \frac{1}{2} \cdot \sup_{f \in X} \inf_{U_n} \text{diam}(F \cap (U_n + f)). \quad (2.5)$$

Both these widths are translation-invariant, and we always have

$$d_{\text{loc}}^n(F) \leq d_{\text{glob}}^n(F). \quad (2.6)$$

If  $F$  is convex and symmetric, then

$$d^n(F) = d_{\text{glob}}^n(F) = d_{\text{loc}}^n(F). \quad (2.7)$$

The widths defined by (2.3)–(2.5) do not increase if  $F$  is replaced by its convex hull. Therefore we can and will assume that  $F$  is convex. In this case the sup over  $X$  in (2.4) and (2.5) clearly can be replaced by a supremum over  $F$ . The global and local Gelfand widths are related to the problem of optimal recovery using nonadaptive or adaptive methods, respectively; see Section 3. Therefore for the adaption problem the following question is interesting. Can the number  $d_{\text{loc}}^n(F)$  be much smaller than  $d_{\text{glob}}^n(F)$ ?

It is useful to study the function

$$f \in F \mapsto \frac{1}{2} \inf_{U_n} \text{diam}((U_n + f) \cap F).$$

A maximum  $f^*$  of that function is called a worst element of  $F$ . If, in addition,

$$d_{\text{glob}}^n(F) = \frac{1}{2} \inf_{U_n} \text{diam}((U_n + f^*) \cap F),$$

then  $f^*$  is called a center of  $F$ . In this case we have

$$d_{\text{loc}}^n(F) = d_{\text{glob}}^n(F).$$

If  $F$  is convex and symmetric, then 0 is a center of  $F$ . This follows from the fact that

$$\text{diam}((U_n + f) \cap F) = \text{diam}((U_n - f) \cap F) \leq \text{diam}(U_n \cap F)$$

for such an  $F$ . Not every convex set has such a center. It is interesting to know whether every convex set  $F$  can be increased slightly such that the bigger set has a center. We will see that this is not the case.

There are many papers and also books on  $n$ -widths. Nonsymmetric sets have rarely been studied so far, however. This seems to be related to the fact that for Kolmogorov widths and global Gelfand widths nonsymmetric sets do not yield very interesting results. By this we mean the following. Let  $F \subset X$  be convex (and nonsymmetric). Then the symmetric set

$$F - F = \{f_1 - f_2 \mid f_i \in F\}$$

is the smallest symmetric set that "contains"  $F$  (more exactly:  $F - F$  contains a translation of  $F$ ). The following result says that the  $n$ -widths of  $F$  and its symmetrization  $F - F$  differ at most by a constant of two.

PROPOSITION 1. *Let  $F \subset X$  be convex. Then*

$$d_n(F) \leq d_n(F - F) \leq 2d_n(F),$$

where  $d_n$  is as given in (2.3), and

$$d_{\text{glob}}^n(F) \leq d^n(F - F) \leq 2d_{\text{glob}}^n(F). \quad (2.8)$$

*Proof.* We only give the proof of (2.8). The inequalities for the Kolmogorov widths are even easier to prove. Assume that  $U_n$  is a closed subspace of  $X$  with codimension  $n$  such that

$$\delta = \sup_{f \in X} \text{diam}((F - F) \cap (U_n + f)).$$

For any fixed  $f_1 \in F$  we get

$$\delta \geq \sup_{f \in X} \text{diam}((F - f_1) \cap (U_n + f)) = \sup_{f \in X} \text{diam}(F \cap (U_n + f)).$$

This proves  $d_{\text{glob}}^n(F) \leq d^n(F-F)$ . Let  $U_n$  be such that

$$\delta = \sup_{f \in X} \text{diam}(F \cap (U_n + f)).$$

Each  $f^* \in (F-F) \cap U_n$  can be written as  $f^* = f_1 - f_2$  with  $f_i \in F$  and also  $f_i \in U_n + f$ . Because of the assumption we obtain  $\|f^*\| \leq \text{diam}(F \cap (U_n + f)) \leq \delta$ . Hence

$$\text{diam}((F-F) \cap U_n) \leq 2\delta.$$

Because  $F-F$  is symmetric we can conclude that

$$\text{diam}((F-F) \cap (U_n + f)) \leq 2\delta$$

for each  $f$  and therefore  $d^n(F-F) = d_{\text{glob}}^n(F-F) \leq 2d_{\text{glob}}^n(F)$ . ■

Such a result does not hold for the local widths. The following example also shows that the local widths can be much smaller than the global widths.

EXAMPLE 1. Let  $X = l_\infty$  and

$$F = \left\{ x \in X \mid x_i \geq 0, \sum_{i=1}^{\infty} x_i \leq 1 \right\}.$$

Then we have

$$d_{\text{glob}}^n(F) = \frac{1}{2}$$

for every  $n \in \mathbb{N}$  and also  $d^n(F-F) = 1$ . For the local widths of  $F$ , however, we obtain

$$\frac{1}{2n+2} \leq d_{\text{loc}}^n(F) \leq \frac{1}{n+1}.$$

*Proof.* Because  $\text{diam}(F) = 1$  we certainly have  $d_{\text{glob}}^n(F) \leq \frac{1}{2}$  and  $d^n(F-F) \leq 1$  for all  $n \in \mathbb{N}$ . One can also prove that  $d^n(F-F) \geq 1$ . Because I did not find a reference for this result, I sketch a proof using results and notation of Section 4. There we prove  $e_{\text{ad}}^n(S|_F) = \frac{1}{2}$ , the statement  $e_{\text{ad}}^n(S|_{F-F}) = e_{\text{non}}^n(S|_{F-F}) = 1$  can be proved analogously or even more easily. Proposition 3 then implies  $d^n(F-F) \geq 1$ . Hence we obtain  $d_{\text{glob}}^n(F) = \frac{1}{2}$  from (2.8).

Now we study the local widths. Consider the set

$$\{x \in l_\infty \mid x_i \in [0, 1/(n+1)] \text{ for } i = 1, \dots, n+1, \text{ and } x_i = 0 \text{ for } i > n+1\} \subset F.$$

This set is (up to translation) an  $(n+1)$ -dimensional ball of diameter  $1/(n+1)$ , and it is well known that the  $n$ -width of such a set is  $1/(2n+2)$ . It then easily follows that

$$d_{\text{loc}}^n(F) \geq \frac{1}{2n+2}.$$

Now fix an  $x \in F$ . There is a partition  $\mathbf{N} = N_1 \cup N_2 \cup \dots \cup N_n$  with the property that for each  $k$  either  $N_k$  contains only one number or

$$L_k(x) := \sum_{i \in N_k} x_i \leq \frac{2}{n+1}.$$

It is clear that the  $N_k$  and the  $L_k$  depend on  $x \in F$ . Assume that  $y \in F$  with  $L_k(y) = L_k(x)$  for all  $k$ . It follows, in particular, that for any  $i$  we have  $y_i = x_i$  or  $y_i \in [0, 2/(n+1)]$ . We obtain

$$\text{diam}(F \cap \{y \mid L_k(y) = L_k(x), k = 1, \dots, n\}) \leq \frac{2}{n+1}$$

and this implies

$$d_{\text{loc}}^n(F) \leq \frac{1}{n+1}. \quad \blacksquare$$

*Remark.* We have seen that the local widths can be much smaller than the global widths. Is there a bound on how much smaller they can be? In other words, is there an inequality of the form

$$d_{\text{glob}}^n(F) \leq c_n \cdot d_{\text{loc}}^n(F) \tag{2.9}$$

with a sequence  $c_n$  that is independent of  $F$ ? If this is the case then of course it would be interesting to know the best inequality of the form (2.9). Actually we conjecture that Example 1 is the most extreme example in the sense that  $d_{\text{glob}}^n(F) = O(n \cdot d_{\text{loc}}^n(F))$  is always true. This is a deep problem which seems to be related to a conjecture of Mityagin and Henkin [21], see also [4], which is still open.

### 3. DIAMETERS OF MAPPINGS AND OPTIMAL RECOVERY

Now we define the widths of  $S|_F$ , where  $S: X \rightarrow Y$  is a continuous linear mapping into a normed space  $Y$ . We use the notation

$$d_n(S|_F) = \inf_{Y_n} \sup_{f \in F} \inf_{g \in Y_n} \|S(f) - g\|, \tag{3.1}$$



where  $Y_n$  runs over all  $n$ -dimensional affine subspaces of  $Y$ , for the Kolmogorov widths. The global Gelfand width (for arbitrary  $F \subset X$ ) is given by

$$d_{\text{glob}}^n(S|_F) = \frac{1}{2} \cdot \inf_{U_n} \sup_{f \in X} \text{diam}(S(F \cap (U_n + f))), \quad (3.2)$$

while a "local" variant is given by

$$d_{\text{loc}}^n(S|_F) = \frac{1}{2} \cdot \sup_{f \in X} \inf_{U_n} \text{diam}(S(F \cap (U_n + f))). \quad (3.3)$$

Here the infimum is taken over closed subspaces  $U_n$  of  $X$  of codimension at most  $n$ . Now an  $f^*$  can be called a worst element if the supremum in (3.3) is attained for  $f^*$ . If, in addition, the local width equals the global width then  $f^*$  is a "center" of  $F$  with respect to  $S$ . This case is similar to the symmetric case insofar as then adaption can help at most by a factor of two; see Proposition 3. It is useful in what follows to define Bernstein widths. The idea of these widths was already used in the proof of Proposition 1 and in Example 1. Here the definition is

$$b_n(S|_F) = \sup\{r \mid S(F) \text{ contains an } (n+1)\text{-dimensional ball with radius } r\}.$$

Observe that in the case  $S = \text{Id}: X \rightarrow X$  we obtain

$$s_n(S|_F) = s_n(F),$$

where  $s_n$  is one of the widths considered here. This means that the diameters of sets are just special cases of this more general notion.

We study the problem of optimal recovery of  $S(f)$  for  $f \in F \subset X$ , if only (adaptive or nonadaptive) information of the form

$$N(f) = (L_1(f), L_2(f), \dots, L_n(f))$$

is available. Each method is of the form  $S_n^{\text{ad}} = \phi \circ N$  with some  $\phi: \mathbf{R}^n \rightarrow Y$  and we want to minimize the maximal error

$$A_{\text{max}}(S_n^{\text{ad}}) = \sup_{f \in F} \|S(f) - S_n^{\text{ad}}(f)\|.$$

In this section we assume that the  $L_i$  are arbitrary linear continuous functionals  $L_i: X \rightarrow \mathbf{R}$ . In the adaptive case the choice of  $L_i$  may depend on  $L_1(f), \dots, L_{i-1}(f)$ . See, for example, Traub, Wasilkowski, Woźniakowski [35] for the exact definitions and known results. If we consider only

methods  $S_n = \phi \circ N$  with a fixed information mapping  $N$ , then we obtain the radius of  $N$  by

$$\text{rad}(N, F) = \inf_{\phi} \Delta_{\max}(\phi \circ N).$$

In connection with Proposition 1 we have the following result.

**PROPOSITION 2.** *Assume that  $N: X \rightarrow \mathbf{R}^n$  is nonadaptive information. Then we have*

$$\text{rad}(N, F) \leq \text{rad}(N, F - F) \leq 4 \text{rad}(N, F).$$

*Proof.* This just follows from

$$\begin{aligned} \text{rad}(N, F) &\leq \text{rad}(N, F - F) \leq \text{diam}(N, F - F) \\ &\leq 2 \text{diam}(N, F) \leq 4 \text{rad}(N, F), \end{aligned}$$

where  $\text{diam}(N, F)$  is defined by

$$\text{diam}(N, F) = \sup_y \text{diam}\{Sf \mid Nf = y\}.$$

The constant 4 is probably not optimal here. ■

In the following we want to compare the numbers

$$e_{\text{non}}^n(S|_F) = \inf \Delta_{\max}(S_n)$$

with the numbers

$$e_{\text{ad}}^n(S|_F) = \inf \Delta_{\max}(S_n^{\text{ad}}),$$

where  $S_n$  runs through all nonadaptive methods and  $S_n^{\text{ad}}$  runs through all adaptive methods using information  $N$  consisting of  $n$  linear functionals. We always assume that  $F$  is convex. First we note a connection between these error bounds and the Gelfand widths. We skip the simple proof.

**PROPOSITION 3.** *Let  $F \subset X$  be a convex set and let  $S: X \rightarrow Y$  be linear and continuous. Then*

$$\frac{1}{2} \cdot d_{\text{glob}}^n(S|_F) \leq e_{\text{non}}^n(S|_F) \leq d_{\text{glob}}^n(S|_F)$$

and

$$\frac{1}{2} \cdot d_{\text{loc}}^n(S|_F) \leq e_{\text{ad}}^n(S|_F).$$

*Remarks.* (a) Assume that  $F$  is convex and symmetric. Then the result is well known. Because of (2.7) we have

$$d_{\text{glob}}^n(S|_F) = d_{\text{loc}}^n(S|_F)$$

in this case and it follows that adaption can help only by a factor of two since

$$e_{\text{non}}^n(S|_F) \leq d_{\text{glob}}^n(S|_F) = d_{\text{loc}}^n(S|_F) \leq 2e_{\text{ad}}^n(S|_F).$$

This is known from [7, 36]. See Kon and Novak [13, 14] and Traub, Wasilkowski, Woźniakowski [35] for further results.

(b) In the next section we will present examples where adaptive methods are much better than nonadaptive ones.

#### 4. ON THE ADAPTION PROBLEM

Our next example shows that adaptive methods may be much better than nonadaptive methods. This example was constructed to demonstrate the superiority of adaptive methods. We do not know, whether it is “worst possible” for nonadaptive methods. We mention that a similar example is contained in Novak [25].

EXAMPLE 2. Let  $X = l_x$  and

$$F = \left\{ x \in X \mid x_i \geq 0, \sum_{i=1}^x x_i \leq 1, x_k \geq x_{2k}, x_k \geq x_{2k+1} \right\}.$$

Let  $e^i$  be the sequence defined by  $e_k^i = \delta_{ik}$ . For  $m \in \mathbb{N}$  we obtain

$$e^i/m \in F - F, \quad i = 1, \dots, 2^{m-1}.$$

Now we use a result of Kashin [11] (see Pinkus [28]) on the Gelfand numbers of the octahedron  $O_n = \{x \in \mathbb{R}^n \mid \sum |x_i| \leq 1\}$  in the  $l_x$ -norm, namely

$$d^n(O_{2n}) \asymp 1/\sqrt{n}.$$

Therefore for  $m \in \mathbb{N}$  we conclude that

$$d^{2^{m-2}}(F - F) \geq \frac{1}{m} \cdot d^{2^{m-2}}(O_{2^{m-1}}) \geq \frac{c'}{2^{m/2}m}.$$

From this we easily derive the lower bound

$$d_{\text{glob}}^n(F) \asymp d^n(F - F) \geq \frac{c}{\sqrt{n \log n}}$$

for the global Gelfand width of  $F$ . For the local widths we clearly have the estimate

$$d_{\text{loc}}^n(F) \leq \frac{1}{2n+2},$$

since  $F$  is contained in the set which was studied in Example 1. Assume now that we want to reconstruct  $x \in F$  in  $l_\infty$ -norm using (adaptively or nonadaptively) linear functionals as information. That is,  $S = \text{Id}$ . For the error of optimal nonadaptive methods we have the lower bound

$$e_{\text{non}}^n(S|_F) = \frac{1}{2} d_{\text{glob}}^n(F) \geq \frac{c}{\sqrt{n \log n}}.$$

Now we describe an adaptive method which is much better. For simplicity we assume here that  $n = 2m - 1$  is odd. By  $\delta_i$  we mean the functional  $\delta_i(x) = x_i$ . First we describe the functionals  $L_i$  which are of the form  $L_i = \delta_{l_i}$ . Take  $L_1 = \delta_1$ . Suppose that  $L_i(x) = x_{l_i}$  are already computed for  $1 \leq i \leq 2k - 1$ . Define

$$J_k = \{j \in \{l_1, \dots, l_{2k-1}\} \mid 2j \notin \{l_1, \dots, l_{2k-1}\}\}$$

and

$$j_k = \min \left\{ j \in J_k \mid x_j = \max_{l \in J_k} x_l \right\}.$$

Take  $L_{2k} = \delta_{2j_k}$  and  $L_{2k+1} = \delta_{2j_k+1}$ , i.e.,  $l_{2k} = 2j_k$  and  $l_{2k+1} = 2j_k + 1$ . We obtain

$$J_{k+1} = J_k \cup \{2j_k, 2j_k + 1\} \setminus \{j_k\}$$

and from  $x_{j_k} \geq x_{2j_k}$  and  $x_{j_k} \geq x_{2j_k+1}$  we conclude that

$$x_{j_{k+1}} \leq x_{j_k}. \quad (4.1)$$

We consider the adaptive information

$$\begin{aligned} N_n(x) &= (L_1(x), \dots, L_n(x)) = (x_{l_1}, \dots, x_{l_n}) \\ &= (x_1, x_2, x_3, x_{2j_2}, x_{2j_2+1}, \dots, x_{2j_{m-1}+1}). \end{aligned}$$

The  $l_i$  are distinct because, aside from  $l_1$ , they come in pairs of the form  $2j_k$  and  $2j_k + 1$ , where  $2j_k$  is always an index not chosen before.

The set  $J_k$  has the following property. Assume that  $l$  is any coordinate not contained in  $\{l_1, l_2, \dots, l_{2k-1}\}$ . Then there is a  $j \in J_k$  with  $x_l \leq x_j$ . The simple proof of this fact would use induction and a partial order on  $\mathbb{N}$ , defined by the dyadic expansion. Consider, for example, the case where  $l = 86$  is not contained in  $\{l_1, l_2, \dots, l_{2k-1}\}$ . Then, depending on  $k$ , exactly one of the indices 1, 2, 5, 10, 21, or 43 is in  $J_k$  and the respective coordinate is at least as large as  $x_l$ . In particular we obtain  $x_l \leq x_{j_k}$  and  $j_k \notin J_{k+1}$ .

Together with (4.1) this means that  $x_l \leq x_j$  for at least  $k$  distinct values of  $j$  and we obtain

$$x_l \leq \frac{1}{k+1}.$$

So for  $n = 2m - 1$  we get  $x_l \leq 1/(m+1) = 2/(n+3)$  for all  $l$  different from the  $x_{l_1}, \dots, x_{l_n}$ . Hence we can reconstruct  $x$  from the information  $N_n$  up to an error of  $1/(2m+2) = 1/(n+3)$ , i.e., we have found a method with

$$\Delta_{\max}(S_n^{\text{ad}}) \leq 1/(n+3).$$

EXAMPLE 1 (continuation). In Proposition 3 we obtained only a one-sided estimate of the error of optimal adaptive methods through the local widths. Can we also prove an upper bound for the error of optimal adaptive methods through the local widths? To answer this question it is enough to study  $S = \text{Id}: F \rightarrow l_\infty$ , where  $F$  is as in Example 1. We claim that

$$e_{\text{ad}}^n(S|_F) = \frac{1}{2}$$

holds for all  $n$ . This means that we have an example with  $e_{\text{ad}}^n(S|_F) \asymp n \cdot d_{\text{loc}}^n(S|_F)$ .

*Proof.* Because  $\text{diam}(F) = 1$  it is enough to show that  $e_{\text{ad}}^n(S|_F) \geq \frac{1}{2}$ . Assume that  $N: F \rightarrow \mathbb{R}^n$  is some adaptive information. We have to prove that

$$\sup_{z \in \mathbb{R}^n} \text{diam}\{x \in F \mid N(x) = z\} = 1.$$

For each  $x \in F$ , let

$$N(x) = (L_1(x), \dots, L_n(x)),$$

where  $L_i$  depends on  $L_1, \dots, L_{i-1}$  and on the values  $L_1(x), \dots, L_{i-1}(x)$ . The  $L_i$  can also be considered as functionals on  $c_0$ , the space of convergent sequences, and therefore are of the form

$$L_i(x) = \sum_{j=1}^{\infty} a_j^i x_j \quad (\text{for } x \in F)$$

with  $(a_j^i)_{j \in \mathbf{N}} \in l_1$ . Now let  $j_1$  be an index such that

$$|a_{j_1}^1| \geq |a_j^1| \quad \text{for all } j.$$

We can assume that the functional  $L_1$  has the form

$$L_1(x) = x_{j_1} + \sum_{j \neq j_1} a_j^1 x_j \quad (4.2)$$

with  $|a_j^1| \leq 1$  for all  $j$ . Observe that (4.2) is true for all  $x \in F := F_0$  because the first functional  $L_1$  is independent of  $x$ . Consider the  $x \in F_0$  for which  $L_1(x) = c_1$  for some small positive  $c_1$  and put

$$F_1 := \{x \in F_0 \mid L_1(x) = c_1\}.$$

For our later analysis it is convenient to assume that

$$c_1 = 3^{n-1} \cdot \delta,$$

where  $\delta > 0$  is sufficiently small that

$$2n \cdot 3^{n-1} \delta < 1. \quad (4.3)$$

In this case the set  $F_1$  is not empty. The second functional  $L_2$  is well-defined on  $F_1$  since it only depends on  $L_1$  and  $c_1$ . We can assume that this linear functional is of the form  $x \mapsto \sum_{j=1}^{\infty} a_j^2 x_j$  with  $a_{j_1}^2 = 0$ . If  $a_{j_1}^2 \neq 0$  then we can consider  $L_2 + \alpha L_1$  instead of the original  $L_2$  without essentially changing the information operator. By multiplying with a constant we now even assume that

$$L_2(x) = x_{j_2} + \sum_{j \notin \{j_1, j_2\}} a_j^2 x_j$$

with  $|a_j^2| \leq 1$  for all  $j$ . Let

$$c_2 = 3^{n-2} \cdot \delta$$

and consider the set

$$F_2 := \{x \in F_1 \mid L_2(x) = c_2\}.$$

The set  $F_2$  is nonempty since there is a (unique)  $x \in F_2$  whose only nonzero coordinates are  $x_{j_1}$  and  $x_{j_2}$ . The functional  $L_3$  is well-defined on  $F_2$  since it only depends on  $L_1, L_2, c_1$ , and  $c_2$ . For general  $k$  we define

$$F_k := \{x \in F_{k-1} \mid L_k(x) = c_k\},$$

where  $L_k$  is the next functional dependent on  $L_1, \dots, L_{k-1}$  and  $c_1, \dots, c_{k-1}$  and  $c_k$  is defined by

$$c_k = 3^{n-k} \cdot \delta. \quad (4.4)$$

The set  $F_k$  is not empty since there is a (unique)  $x \in F_k$  whose only nonzero coordinates are  $x_{j_1}, \dots, x_{j_k}$ .

We can and will assume that each  $L_k$  is of the form

$$L_k(x) = x_{j_k} + \sum_{j \in \{j_1, \dots, j_k\}} a_j^k x_j$$

with distinct  $j_k$  and  $|a_j^k| \leq 1$  for all  $j$ .

Let  $k_n$  be an index for which the coordinates  $a_{k_n}^i$  are small for all  $i = 1, \dots, n$ , i.e.,

$$|a_{k_n}^i| < \delta \quad \text{for } i = 1, \dots, n. \quad (4.5)$$

The sequences  $(a^i)_j$  are in  $c_0$  and therefore such an index  $k_n$  clearly exists. We consider a sequence  $x$  with nonzero coordinates  $x_{j_1}, \dots, x_{j_n}$  and  $x_{k_n} = c \in [0, 1]$ ; i.e., we assume that  $x_i = 0$  for all other  $i$ . We also assume that

$$L_i(x) = c_i \quad (i = 1, \dots, n).$$

It is easy to see that the sequence  $x \in l_x$  is uniquely defined by these requirements for each given  $c$ . The nonzero coordinates of the sequence  $x$  are given by the linear system

$$\begin{aligned} x_{j_n} &= -a_{k_n}^n c + c_n \\ a_{j_n}^{n-1} x_{j_n} + x_{j_{n-1}} &= -a_{k_n}^{n-1} c + c_{n-1} \\ &\dots \\ a_{j_n}^1 x_{j_n} + a_{j_{n-1}}^1 x_{j_{n-1}} + \dots + a_{j_2}^1 x_{j_2} + x_{j_1} &= -a_{k_n}^1 c + c_1. \end{aligned}$$

We also have  $x_{k_n} = c$  while all other coordinates are zero. By (4.5) and  $0 \leq c \leq 1$  we know that all terms of the form  $|a_{k_n}^i c|$  are bounded by  $\delta$ . Using (4.4), for the unique solution of the system we get the estimates

$$\begin{aligned}
0 < x_{j_n} < 2\delta \\
0 < x_{j_{n-1}} < 6\delta \\
& \dots \\
0 < x_{j_1} < 2 \cdot 3^{n-1} \delta.
\end{aligned}$$

We obtain

$$\sum_{i=1}^n x_{j_i} < 2n \cdot 3^{n-1} \delta$$

and conclude from (4.3) that  $x \in F_n$  for all  $c = x_{k_n}$  with

$$0 \leq c \leq 1 - 2n \cdot 3^{n-1} \delta.$$

Thus  $\sup_{z \in \mathbf{R}^n} \text{diam}\{x \in F \mid N(x) = z\} \geq 1 - 2n \cdot 3^{n-1} \delta$  for each  $\delta > 0$  as in (4.3). Letting  $\delta \rightarrow 0$ , we obtain our result. ■

## 5. THE CASE OF RESTRICTED INFORMATION

In many practical cases,  $X$  is a Banach space of functions and only certain linear functionals are available as information. Here we only consider the case where all functionals

$$\delta_x : f \in X \mapsto f(x) \in \mathbf{R}$$

are continuous and form the set of available functionals; i.e., each  $L_i$  is of the form  $L_i = \delta_{x_i}$ . Hence we study methods of the form

$$S_n^{(\text{ad})}(f) = \phi(f(t_1), \dots, f(t_n)) \quad (5.1)$$

and define the error bounds

$$\tilde{e}_{\text{non}}^n(S|_F) = \inf_{S_n} \mathcal{A}_{\max}(S_n)$$

and

$$\tilde{e}_{\text{ad}}^n(S|_F) = \inf_{S_n^{\text{ad}}} \mathcal{A}_{\max}(S_n^{\text{ad}}),$$

where the infimum runs through all nonadaptive or adaptive methods of the form (5.1), respectively. Here we only give some examples and therefore we do not define any new  $n$ -widths with a restriction to the type of admissible subspaces. The numbers  $\tilde{e}_{\text{non}}^n(S|_F)$  can be estimated from



above by the Kolmogorov widths  $d_n(F)$ , even if  $S$  is nonlinear; see Novak [22, 23]. The estimate for the case of a linear functional was also found by Belinskii [5].

The following example shows that adaptive methods may be exponentially better than nonadaptive ones.

EXAMPLE 3. Again we consider  $S = \text{Id} : F \rightarrow l_\infty$  with

$$F = \left\{ x \in X \mid x_i \geq 0, \sum_{i=1}^{\infty} x_i \leq 1, x_k \geq x_{2k}, x_k \geq x_{2k+1} \right\},$$

but now we only allow methods of the form (5.1). We already proved that

$$\tilde{e}_{\text{ad}}^n(S|_F) \leq \frac{1}{n+3}$$

and it is not difficult to see that the nonadaptive information

$$N_n(x) = (x_1, x_2, \dots, x_n)$$

is optimal among all nonadaptive information operators. It follows, in particular, that

$$\tilde{e}_{\text{non}}^n(S|_F) \asymp \frac{1}{\log(n+1)}.$$

EXAMPLE 4. We mention the following example from Korneichuk [16], which is closer to the standard classes of approximation theory. Consider the reconstruction problem  $S = \text{Id} : F \rightarrow L_\infty([0, 1])$  for

$$F = \{ f : [0, 1] \rightarrow [0, 1] \mid f \text{ monotone and } |f(x) - f(y)| \leq |x - y|^\alpha \}$$

with  $0 < \alpha < 1$ . This problem can be solved adaptively using the bisection method, while nonadaptive methods are worse. Korneichuk [16] proved, more exactly, that

$$\tilde{e}_{\text{non}}^n(S|_F) \asymp n^{-\alpha} \quad \text{while} \quad \tilde{e}_{\text{ad}}^n(S|_F) \asymp n^{-1} \log n.$$

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